

## Error Correcting Codes and Secret Sharing Intro

**Secret Sharing:** We make use of the fact that there is a unique polynomial of degree  $d$  passing through a given set of  $d + 1$  points. This means that if we require  $k$  people to come together in order to find a secret, we should use a polynomial of degree  $k - 1$ , and give each person one point. There are more complicated schemes if there are more conditions, but they all use the same concept.

**Erasure errors:** A packet/point  $(x, P(x))$  is *lost* in the communication channel.

$$\begin{array}{|c|c|c|c|c|} \hline 1 & 5 & 3 & 4 & 3 \\ \hline \end{array} \longrightarrow \begin{array}{|c|c|c|c|c|} \hline 1 & 5 & & 4 & 3 \\ \hline \end{array}$$

**Protection:** ( $n$  packets,  $k$  errors) interpolate polynomial through the message points, send  $n + k$  packets on the polynomial

**General errors:** A packet is *modified* in the communication channel.

$$\begin{array}{|c|c|c|c|} \hline 1 & 5 & 3 & 4 \\ \hline \end{array} \longrightarrow \begin{array}{|c|c|c|c|} \hline 1 & 5 & 6 & 4 \\ \hline \end{array}$$

**Protection:** ( $n$  packets,  $k$  errors) interpolate polynomial through the message points, send  $n + 2k$  packets on the polynomial

### Berlekamp–Welch Algorithm:

Variables: sent message packets  $m_i$ , received packets  $r_i$ , error locations  $e_i$

Polynomials:

- $P(x)$ : original polynomial through message (this is what we want)
- $E(x) = (x - e_1)(x - e_2) \cdots (x - e_k)$ : error locator polynomial
- $Q(x) = P(x)E(x)$ , or  $Q(i) = P(i)E(i) = r_i E(i)$  for all  $i$

$Q(x)$  and  $E(x)$  are unknown, but we can solve for them using a system of equations.

## 1 Berlekamp-Welch Warm Up

Note 8  
Note 9

Let  $P(i)$ , a polynomial applied to the input  $i$ , be the original encoded polynomial before sent, and let  $r_i$  be the received info for the input  $i$  which may or may not be corrupted.

(a) If you want to send a length- $n$  message, what should the degree of  $P(x)$  be? Why?

(b) When does  $r_i = P(i)$ ? When does  $r_i$  not equal  $P(i)$ ?

- (c) If there are at most  $k$  erasure errors, how many packets should you send? If there are at most  $k$  general errors, how many packets should you send? (We will see the reason for this later.) Now we will only consider general errors.
- (d) What do the roots of the error polynomial  $E(x)$  represent? Does the receiver know the roots of  $E(x)$ ? If there are at most  $k$  errors, what is the maximum degree of  $E(x)$ ? Using the information about the degree of  $P(x)$  and  $E(x)$ , what is the degree of  $Q(x) = P(x)E(x)$ ?
- (e) Why is the equation  $Q(i) = P(i)E(i) = r_i E(i)$  always true? (Consider what happens when  $P(i) = r_i$ , and what happens when  $P(i)$  does not equal  $r_i$ .)
- (f) In the polynomials  $Q(x)$  and  $E(x)$ , how many total unknown coefficients are there? (These are the variables you must solve for. Think about the degree of the polynomials.) When you receive packets, how many equations do you have? Do you have enough equations to solve for all of the unknowns? (Think about the answer to the earlier question - does it make sense now why we send as many packets as we do?)
- (g) If you have  $Q(x)$  and  $E(x)$ , how does one recover  $P(x)$ ? If you know  $P(x)$ , how can you recover the original message?

## 2 Berlekamp-Welch Algorithm

Note 8  
Note 9

In this question we will send the message  $(m_0, m_1, m_2) = (1, 1, 4)$  of length  $n = 3$ . We will use an error-correcting code for  $k = 1$  general error, doing arithmetic over  $\text{GF}(5)$ .

- (a) Construct a polynomial  $P(x) \pmod{5}$  of degree at most 2, so that

$$P(0) = 1, \quad P(1) = 1, \quad P(2) = 4.$$

What is the message  $(c_0, c_1, c_2, c_3, c_4)$  that is sent?

- (b) Suppose you receive the message  $(0, 1, 4, 0, 4)$  and know that one packet was corrupted. Set up the system of linear equations in the Berlekamp-Welch algorithm to find  $Q(x)$  and  $E(x)$ .

- (c) Assume that after solving the equations in part (b) we get  $Q(x) = 4x^3 + x^2 + x$  and  $E(x) = x$ . Show how to recover the original message from  $Q$  and  $E$ .

